

# Algebraic Numbers in Isabelle/HOL<sup>1</sup>

#### René Thiemann and Akihisa Yamada



### Overview

#### Motivation

#### • Real Algebraic Numbers

- Well-Definedness
- Calculating Real Roots of Rational Polynomial
- Factorizing Rational Polynomial
- Arithmetic on Real Algebraic Numbers

#### • Complex Algebraic Numbers

## Certify Complexity of Matrix Interpretations

• given: automatically generated complexity proof for program

 $\chi_A = (x - 1) \cdot (-39 + 360x - 832x^2 + 512x^3)$ 

- criterions
  - polynomial complexity if norms of all complex roots of  $\chi_A\leqslant 1$
  - degree d in  $\mathcal{O}(n^d)$ : more calculations with complex roots of  $\chi_A$

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  - degree d in  $\mathcal{O}(n^d)$ : more calculations with complex roots of  $\chi_A$
- problem: certifier crashed as numbers got too complicated

## Closed Form for Cubic Polynomials

• 
$$-39 + 360x - 832x^2 + 512x^3 = 0$$
 iff  
•  $x = \frac{1}{24} \left( 13 + \frac{34}{\sqrt[3]{91+9i\sqrt{383}}} + \sqrt[3]{91+9i\sqrt{383}} \right),$   
•  $x = \frac{13}{24} - \frac{17(1+i\sqrt{3})}{24\sqrt[3]{91+9i\sqrt{383}}} - \frac{1}{48} (1-i\sqrt{3}) \sqrt[3]{91+9i\sqrt{383}},$  or  
•  $x = \frac{13}{24} - \frac{17(1-i\sqrt{3})}{24\sqrt[3]{91+9i\sqrt{383}}} - \frac{1}{48} (1+i\sqrt{3}) \sqrt[3]{91+9i\sqrt{383}}$ 

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$$norm(x_j) = \sqrt{Re(x_j)^2 + Im(x_j)^2} \leqslant 1$$

for all  $j \in \{1, 2, 3\}$ 

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• problem: no closed form for roots of polynomials of degree 5 and higher

## Algebraic Numbers

• number  $x \in \mathbb{R} \cup \mathbb{C}$  is algebraic iff it is root of non-zero rational polynomial

• 
$$x_1 =$$
 "root #1 of  $-39 + 360x - 832x^2 + 512x^3$ "

• 
$$x_2 =$$
 "root #2 of  $-39 + 360x - 832x^2 + 512x^3$ "

•  $x_3 =$  "root #3 of  $-39 + 360x - 832x^2 + 512x^3$ "



Figure:  $-39 + 360x - 832x^2 + 512x^3$ 

• well-definedness

is there a "root #3 of  $-23 + x - 5x^2 + x^3$ "

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• arithmetic

calculate a polynomial representing "root #2 of  $-2 + x^{2"}$  + "root #1 of  $-3 + x^{2"}$ "

 well-definedness Sturm's method is there a "root #3 of  $-23 + x - 5x^2 + x^{3}$ " factorization of rational polynomials representation is there a simpler representation of "root #3 of  $-108 - 72x + 108x^{2} + 84x^{3} - 27x^{4} - 32x^{5} - 2x^{6} + 4x^{7} + x^{8''}$ Sturm's method comparisons is "root #3 of  $-39 + 360x - 832x^2 + 512x^3$ " smaller than 1 arithmetic matrices, determinants, resultants, .... calculate a polynomial representing "root #2 of  $-2 + x^{2}$ " + "root #1 of  $-3 + x^{2}$ "

## Main Result: Formalization of Algebraic Numbers

- common properties on algebraic numbers ( $\mathbb{R}$  and  $\mathbb{C}$ )
- executable real algebraic numbers
- executable complex algebraic numbers indirectly
- easy to use via data-refinement for  ${\mathbb R}$  and  ${\mathbb C}$

 $\lfloor norm(\sqrt[3]{2}+3+2i) \cdot 100 \rfloor \stackrel{evaluate}{\hookrightarrow} 216$ 

• applicable inside (eval) and outside Isabelle (export-code)

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  - oracle (MetiTarski) performs computations
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  - oracle (MetiTarski) performs computations
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- experimental comparison (examples of Li and Paulson)

MetiTarski	1.83 seconds	(@ 2.66 Ghz)
+ validation of Li and Paulson	4.16 seconds	(@ 2.66 Ghz)
Our generated Haskell code	0.03 seconds	(@ 3.5 Ghz)

#### Well-Definedness



- input
  - polynomial over  $\mathbb R$
  - interval (  $[2,5],\,(-\pi,7],\,(-\infty,3),$  or  $(-\infty,+\infty)$  )
- output: count-roots p itval
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number of distinct real roots of polynomial over  $\mathbb Q$  in interval over  $\mathbb Q$ 

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- precompute first phase of Sturm's method
- $\Rightarrow$  query many intervals for same polynomial more efficiently

## Applying Sturm's Method for Well-Definedness



count-roots  $(-25 + 155x - 304x^2 + 192x^3)$   $(-\infty, +\infty) = 2$ 

## Representation of Real Algebraic Numbers – Part 1

- quadruple
  - polynomial over Q: p
  - left and right interval bound  $\in \mathbb{Q}$ : *I* and *r*
  - precomputation of Sturm for p: root-info
  - (flag for factorization information on p)

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  - $p \neq 0$
  - root-info = count-roots p
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  - . . .

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- five invariants
  - $p \neq 0$
  - root-info = count-roots p
  - root-info [*I*, *r*] = 1
  - ...
- invariants are ensured via a subtype (typedef)
  - new type of quadruples satisfying invariants
  - lift-definition and transfer for function definitions and proofs

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- 3. determine root-info<sub>i</sub> for  $p_i$  and initial bounds  $l_i$ ,  $r_i$  such that

 $\operatorname{root-info}_i (-\infty, +\infty) = \operatorname{root-info}_i [I_i, r_i]$ 

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root-info<sub>i</sub>  $(-\infty, +\infty)$  = root-info<sub>i</sub>  $[I_i, r_i]$ 

perform bisection on [*l<sub>i</sub>*, *r<sub>i</sub>*] to obtain intervals which all contain a single root of *p<sub>i</sub>*

- input:  $\chi_A$
- computation:
  - $39 399x + 1192x^2 1344x^3 + 512x^4$

output:

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- computation:
  - $(x-1) \cdot (-39 + 360x 832x^2 + 512x^3)$

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• output: 
$$(x - 1, [1, 1])$$

- initial bounds:  $6 = degree \ p_i \cdot max\{ [|c_i|], c_i \ coefficient \ of \ p_i \}$
- computation:
  - $(x-1) \cdot (-39 + 360x 832x^2 + 512x^3)$
  - $todo = \{[-6, 6]\}$  $\begin{array}{c} 40 \\ 20 \\ -20$

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1.2

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,  $(p_i, [0, \frac{3}{8}])$ 

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- does not terminate (e.g., pass  $\lambda x.2$ )
- definition as partial-function to obtain code-equations
- soundness via well-founded induction;
   order based on minimal distance between roots of p<sub>i</sub>

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time/degree of representing polynomials for $\sum_{i=1} \sqrt{r}$			
factorization	<i>n</i> = 8	<i>n</i> = 9	n = 10
none square-free full	2m11s/256 2m14s/256 0.35s/16	22m19s/512 15m31s/384 0.35s/16	12h19m/1024 9h31m/768 0.59s/16

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• time / പ c vonvocanting not .....  $\sum n$ *[*.

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invoke oracle on each factor

$$-12 - 8x + 4x^{2} + 4x^{3} + x^{4} \stackrel{oracle}{=} (-2 + x^{2}) \cdot (6 + 4x + x^{2})$$

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check equality at runtime, not irreducibility

# Simplification of Representation



# Comparison of Real Algebraic Numbers

1. decide whether (p, [l, r]) and (q, [l', r']) encode same number

THE root of p in [l, r] = THE root of q in [l', r'] $\iff$  gcd p q has real root in  $[l, r] \cap [l', r']$ 

 if not, tighten bounds of both numbers via bisection until intervals are disjoint (bisection algorithm again defined by partial-function)

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  - *n*√·: *p*(*x<sup>n</sup>*)
  - addition: resultant(p(x y), q(y))

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- 4. optimize representation

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- factorizations in between to simplify representations
- efficient (not optimal) computation of resultants and GCDs
- tuned algorithms on polynomials
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```
datatype real_alg_3 =
  Rational rat
  Irrational "quadruple with invariants"
```

typedef real\_alg\_4 = "real\_alg\_3 with invariant"

definition real\_of\_4 :: real\_alg\_4 => real

quotient\_type real\_alg =
real\_alg\_4 / "% x y. real\_of\_4 x = real\_of\_4 y"

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- missing: equivalent of Sturm's method for  $\ensuremath{\mathbb{C}}$
- take existing Cartesian representation: (Re(x), Im(x))
- only new difficulty: find complex roots, e.g., of  $p = 1 + x + x^3$
- basic idea:
  - complex roots of rational polynomials come in complex conjugate pairs
  - $\Rightarrow$  if x is root of p then  $\bar{x}$  is root of p

- $\bullet\,$  missing: equivalent of Sturm's method for  $\mathbb C$
- take existing Cartesian representation: (Re(x), Im(x))
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  - $\Rightarrow$  compute all these roots

$$Re(x) = \left(SOME \ x. \ 8 - 24x - 88x^3 + 96x^4 + 288x^5 + 384x^6 + 768x^7 + 512x^9 = 0\right)$$
$$Im(x) = \left(SOME \ x. \ \frac{961}{4096}x^2 - \frac{279}{512}x^4 + \frac{453}{256}x^6 - \frac{85}{32}x^8 + \frac{27}{8}x^{10} - 3x^{12} + x^{14} = 0\right)$$
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and filter: for all candidates z = Re(x) + Im(x)i, test p(z) = 0

## Summary

- formalization of real and complex algebraic numbers, implementing ℝ and ℂ
  +, -, ×, /, <sup>n</sup>√·, [·], =, <, *i*, *Re*, *Im*, show
- factorization algorithms for rational polynomials over  $\mathbb{Q},\,\mathbb{R},\,\mathbb{C}$
- heavily relying on Sturm's method and matrix library
- based on factorization oracle (S. Joosten @ Isabelle workshop)
- $\sim$  20.000 loc, available in archive of formal proofs